# Higher order derivatives 

OR CUBE, OR ALGEBRAIC, OR INTEGRAL

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kodeketa eta kriptografiaren egunak Hendaye, le 14 avril 2022


## Outline

REPRESENTATION

Degree

DIVISION PROPERTY

AtTACK STRATEGIES

RANDOM DIRECTIONS

## The Algebraic Normal Form

$f: \mathbb{F}_{2}^{n} \rightarrow \mathbb{F}_{2}$ Then $f$ can be uniquely represented as an element of

$$
\mathbb{F}_{2}\left[x_{0}, \ldots, x_{n-1}\right] /\left(x_{0}^{2}-x_{0}, \ldots, x_{n-1}^{2}-x_{n-1}\right)
$$

That is a sum of monomials, i.e. for some $u \in \mathbb{F}_{2}^{n}$

$$
x^{u}=\prod_{i=0}^{n-1} x_{i}^{u_{i}}
$$

Example : $x_{0} x_{2} x_{3}=x^{10110000}$

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$$
f\left(x_{0}, \ldots, x_{n-1}\right)=\bigoplus_{u \in \mathbb{F}_{2}^{n}} c_{u} x^{u}
$$

with $c_{u} \in \mathbb{F}_{2}$.

## Truth table and Monomials

$$
\begin{gathered}
f\left(x_{0}, \ldots, x_{n-1}\right)=\bigoplus_{u \in \mathbb{F}_{2}^{n}} c_{u} x^{u} \\
f(a)=\bigoplus_{u<a} c_{u} \text { and } c_{u}=\bigoplus_{a<u} f(a)
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Where $a \prec u$ iff $\operatorname{supp}(a) \subset \operatorname{supp}(u)$ Also

$$
\operatorname{wt}(u)=\# \operatorname{supp}(u)
$$

## Functions

A function from $\mathbb{F}_{2}^{n}$ to $\mathbb{F}_{2}^{m}$ is represented as a collection of $m$ boolean functions, called component functions.

- For permutations, the monomial $x_{0} x_{1} \cdots x_{n-1}$ never appears
- A random function has its monomials appearing each with probability $1 / 2$ in each component function.


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## Higher-order differential attacks [Lai 94]

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DEFINITION (MULTIVARIATE DEGREE)

$$
d=\operatorname{deg}(f)=\max \left\{\operatorname{wt}(u), c_{u}=1\right\}
$$

## Higher-order differential attacks [Lai 94]

$$
f\left(x_{0}, \ldots, x_{n-1}\right)=\bigoplus_{u \in \mathbb{F}_{2}^{\prime}} c_{u} x^{u}
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## DEFINITION (MULTIVARIATE DEGREE)

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$$

Distinguisher :
For all linear space $V$, with $\operatorname{dim}(V) \geq d+1$,

$$
g: x \mapsto \sum_{v \in V} f(x+v)
$$

is constant to zero.

## DEGREE EVALUATION

At different level :

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- Better upper bound when the structure is specific [CGGLRS, 2020]


## Going Further

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A component of $E_{k}(x)$ can be represented as

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\sum_{u \in \mathbb{F}_{2}^{n}} c_{u}(k) x^{u}
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and assume that for any $v \succ u, c_{v}(k)=0$.
Then what is the possible degree of $E_{k}(x)$ in $x$ ?

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f(x)=x_{0} x_{1} x_{2} x_{3} x_{4} x_{6}+x_{3}+x_{4} x_{5}+x_{6}
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Upper bound is not enough : lower bound [HLLT 2020]

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## Specific Higher order derivative

Assume that for any $w \succ u, c_{w}(k)=0$, then let $V=\{v, v \prec u\}$. Then for any $x$,

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> Division Property [Todo 2015]

## Techniques

- Using the representation of the Sbox and the linear layer, this division property can be used for iterated construction
- Mixed Integer Linear Programming
- Lower bound the degree
- Monomial prediction, monomial trails


## Problems

- Easy for one monomial, not easy for all...


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- Not linearly equivalent [LDF, 2020]


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## Block ciphers



## Block ciphers



- Proofs of modes, wrt indistinguishability
- Same reasoning for permutation-based constructions.


## BLOCK CIPHERS



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## TAKING THE MODE INTO ACCOUNT

Pyjamask-96

- Distinguisher integral over $10+1$ rounds
- 3 rounds in the backward direction (monomial count)


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Considering the data complexity...

## ON DUPLEX OR STREAM CIPHERS



## ON DUPLEX OR STREAM CIPHERS



What can you do?
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## RANDOM DIRECTIONS

- Representation of polynomials?
- Given a polynomial, find a (non) linear transofrmation that would become an affine space after application?
- Provide a way to state "every $c_{u}(k)$ is complicated enough"?
- Criteria that would be equivalent under the representation of the transformation?


## TAKE AWAY

I'm sure there is a monomial missing somewhere!

